



# Demand-Aware Route Planning for Shared Mobility Services

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# Outline

- **Motivation**
- Problem Formulation
- Algorithms
- Evaluations



# Background

Development of *shared* mobility.

- Food delivery
- Ridesharing
- Crowdsourced parcel delivery



# Online platforms for shared mobility

- Cainiao
- Meituan
- Uber
- Didi
- ...



# Route Planning for Shared Mobility

- A **large** amount of **dynamically** arriving requests
- A **large** amount of workers





# Route Planning for Shared Mobility



- A **large** amount of possible **route** allowing share
- **Limited** response **time**

# Route Planning for Shared Mobility

- Large amount of **dynamically** arriving requests
- Large amount of workers
- Large amount of possible routes allowing share
- Limited response time

effective / efficient  
**route planning** strategy

# Keep the Balance of Demand-Supply

Great profit **loss** from *unmatched* distribution of *demand* and *supply*

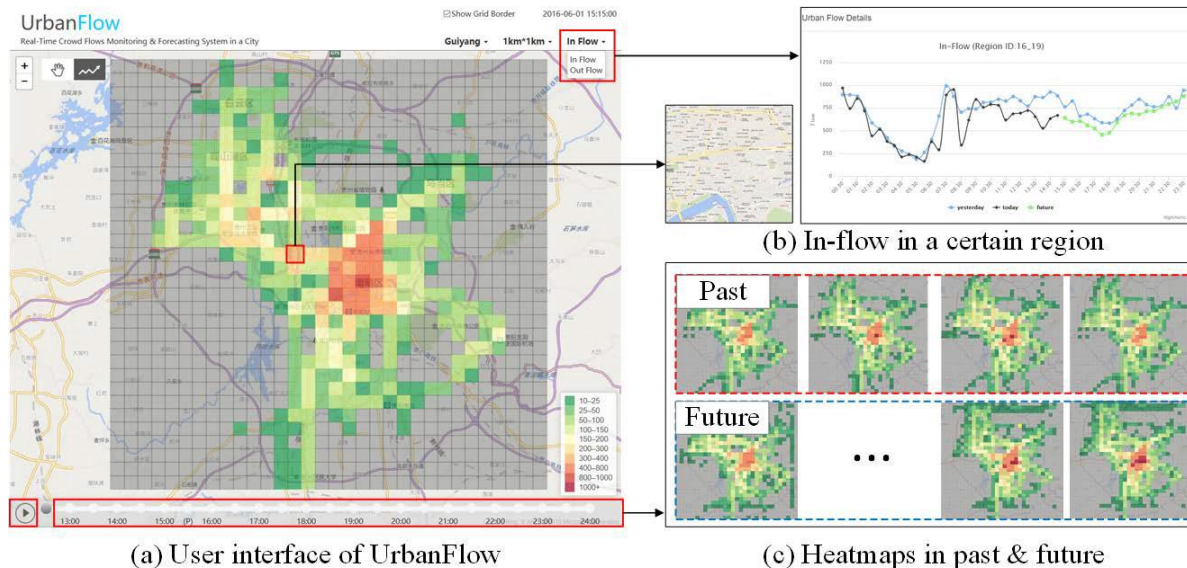
- **Rush hour** (ridesharing)
  - Morning: rural areas → center of the city
  - Evening: center of the city → rural areas
- **Lunch and supper time** (food delivery)
  - Tons of orders sent to business central





# Prediction of Demand

Derive demands in different areas/timesteps accurately using **spatiotemporal** prediction. (e.g. DeepST\*)



(b) In-flow in a certain region

(c) Heatmaps in past & future

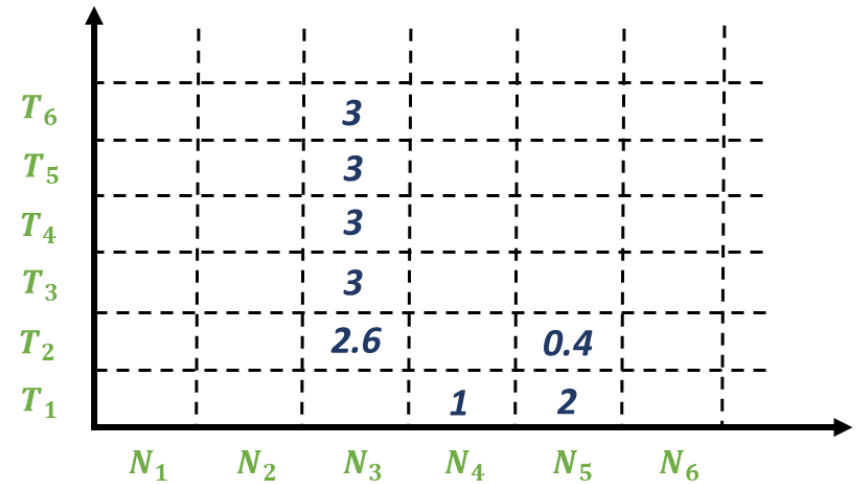
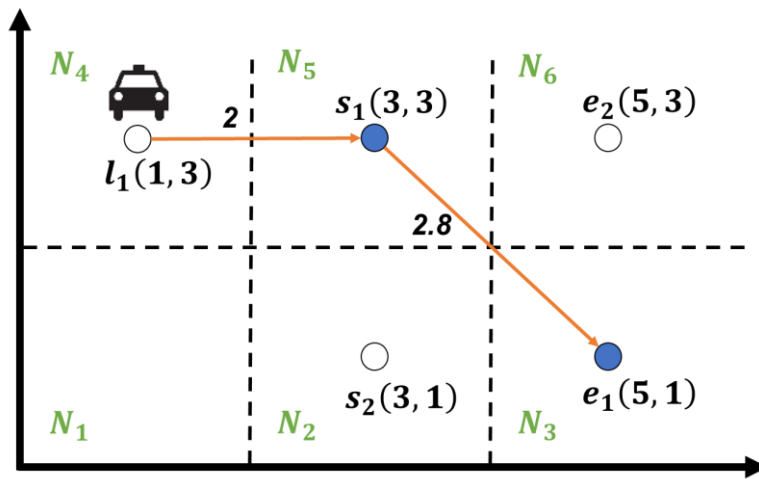
\*Figure is Provided by Junbo Zhang

How to use it to benefit route planning for shared mobility?

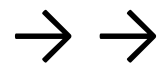


# Supply Organization

In each **area** and **time span**, a larger total **time duration** of workers leads to a higher probability to serve a request.



Different **route planning** strategy

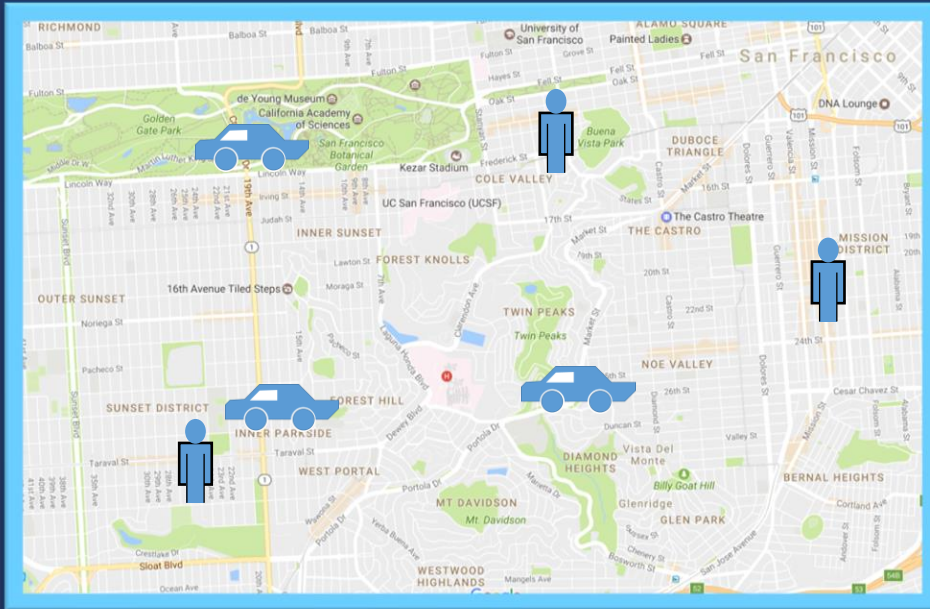


Different **distribution of supply**

Organize **supply** according to **demand** to maximize profit.



# Motivation



Design an algorithm to **improve the effectiveness** of *route planning* for shared mobility through:

Evaluating the effect of **supply** during route planning based on **demand**.

The **overall profit** of the platform is **improved**.



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# Workers and Requests

## Workers

$$W = \{w_1, w_2, \dots, w_n\}$$

$w_i$

current location  $l_i$

capacity  $a_i$

## Requests

$$R = \{r_1, r_2, \dots, r_m\}$$

$r_j$

start/end loc.  $s_j/e_j$

release time  $tr_j$

deadline  $td_j$

rejection penalty  $p_j$

capacity  $a_j$



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# Workers and Requests

Workers

$$W = \{w_1, w_2, \dots, w_n\}$$

Requests

$$R = \{r_1, r_2, \dots, r_m\}$$

Route  $S_i$ , a sequence of  $s_j/e_j$

$w_i$

Planning routes to serve requests.

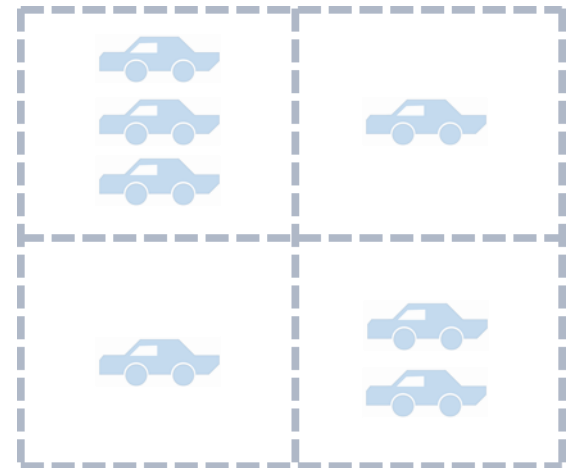
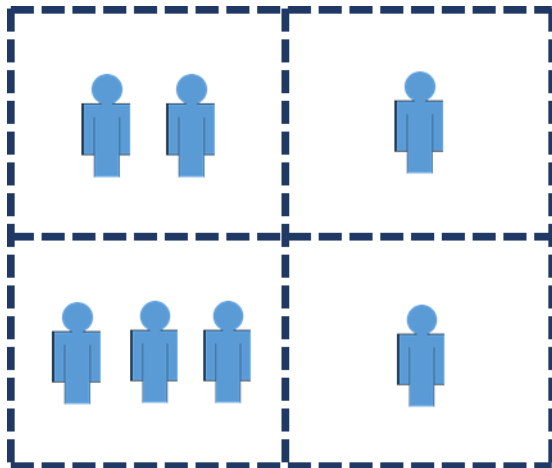
$r_j$

Existing work\* handles the **distance** related cost.

# Quantize Demand-Supply Balance

## 1. Demand number map (DN)

- **Number of requests** in each time span and area
- Predicted using deep learning model\*



# Quantize Demand-Supply Balance

1. Demand number map (DN)

2. Supply number map (SN)

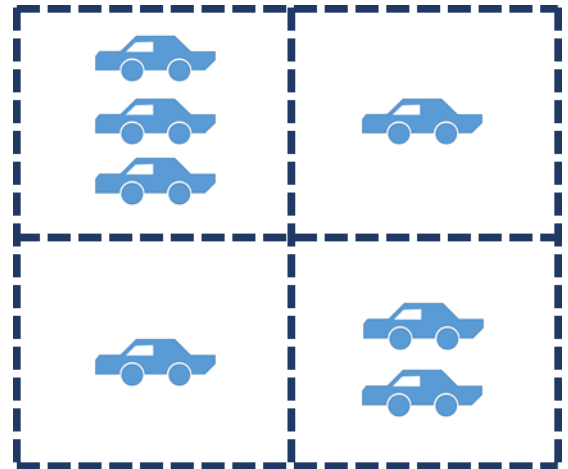
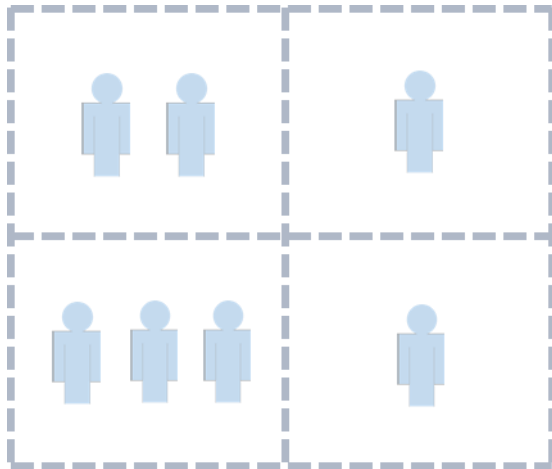
- **Number of workers** in each time span and area
- Updated according to **route planning**

Route plan affects  
Supply Contribution

Supply number  $SN$

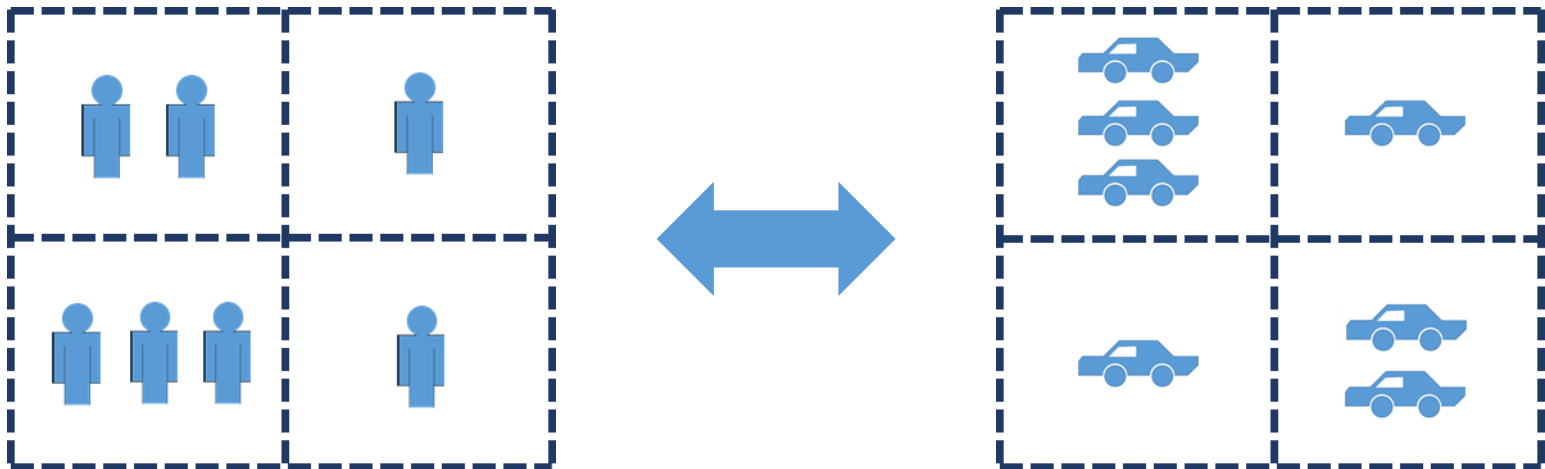
$$SN(N_x, T_y) = \sum_{w_i \in W} \sum_{k=1}^{|S_{w_i}|-1}$$

$SC_{xy}(l_k, l_{k+1}, arr_i[l_k])$



# Quantize Demand-Supply Balance

1. Demand number map (DN)
2. Supply number map (SN)
3. Demand-Supply Balance Score (DSB)
  - Each route plan affects future supply and balance
  - Statistically analyze the expected **profit** of the balance



# Quantize Demand-Supply Balance

1. Demand number map (DN)

2. Supply number map (SN)

3. Demand-Supply Balance Score (DSB)

- Each route plan affects future supply and balance
- Statistically analyze the expected **profit** of the balance

Local Balance  $LB(\lambda, sn)$  =  $E(Y)$  =  $\sum_{k=0}^{\lfloor sn \rfloor - 1} k \frac{\lambda^k}{k!} e^{-\lambda} + \sum_{k=\lfloor sn \rfloor}^{\infty} \lfloor sn \rfloor \frac{\lambda^k}{k!} e^{-\lambda}$  (3)

(Expected # of matchings)

Local demand ( $\lambda$ )  
and supply ( $sn$ )

$$= \sum_{k=0}^{\lfloor sn \rfloor - 1} (k - \lfloor sn \rfloor) \frac{\lambda^k}{k!} e^{-\lambda} + \sum_{k=0}^{\infty} \lfloor sn \rfloor \frac{\lambda^k}{k!} e^{-\lambda}$$

$$= \sum_{k=0}^{\lfloor sn \rfloor - 1} (k - \lfloor sn \rfloor) \frac{\lambda^k}{k!} e^{-\lambda} + \lfloor sn \rfloor$$

Analyzed based on  
**Poisson distribution**

$$DSB = \sum_{x=1}^{|\mathcal{N}|} \sum_{y=1}^{|\mathcal{T}|} \beta_y (sn'_{xy} - sn_{xy}) \cdot \Delta_{LB}(\lambda_{xy}, sn_{xy})$$

Change of  $LB$

# Demand-Aware Route Planning (DARP) Problem

Given a set of workers  $W$ , a set of requests  $R$ , a demand number map  $DN$ , the DARP Problem is to find the sets of routes  $S$  for all the workers to minimize **Demand-Aware Cost (DAC)**:

$$DAC(W, R, DN) = \underbrace{\alpha \sum_{w_i \in W} D(S_{w_i})}_{\text{Cost from workers' moving distances}} - \underbrace{DSB(\beta, S, DN)}_{\text{Cost from Demand-Supply Balance}} + \underbrace{\sum_{r_j \in \bar{R}} p_j}_{\text{Cost from rejection}}$$

Such that:

1. at any time the total capacity of requests of any worker should not exceed its **capacity**  $a_i$ ;
2. each request meets its **deadline**;
3. an assigned request cannot be assigned to another; a rejected request cannot be revoked.





# Demand-Aware Route Planning (DARP) Problem

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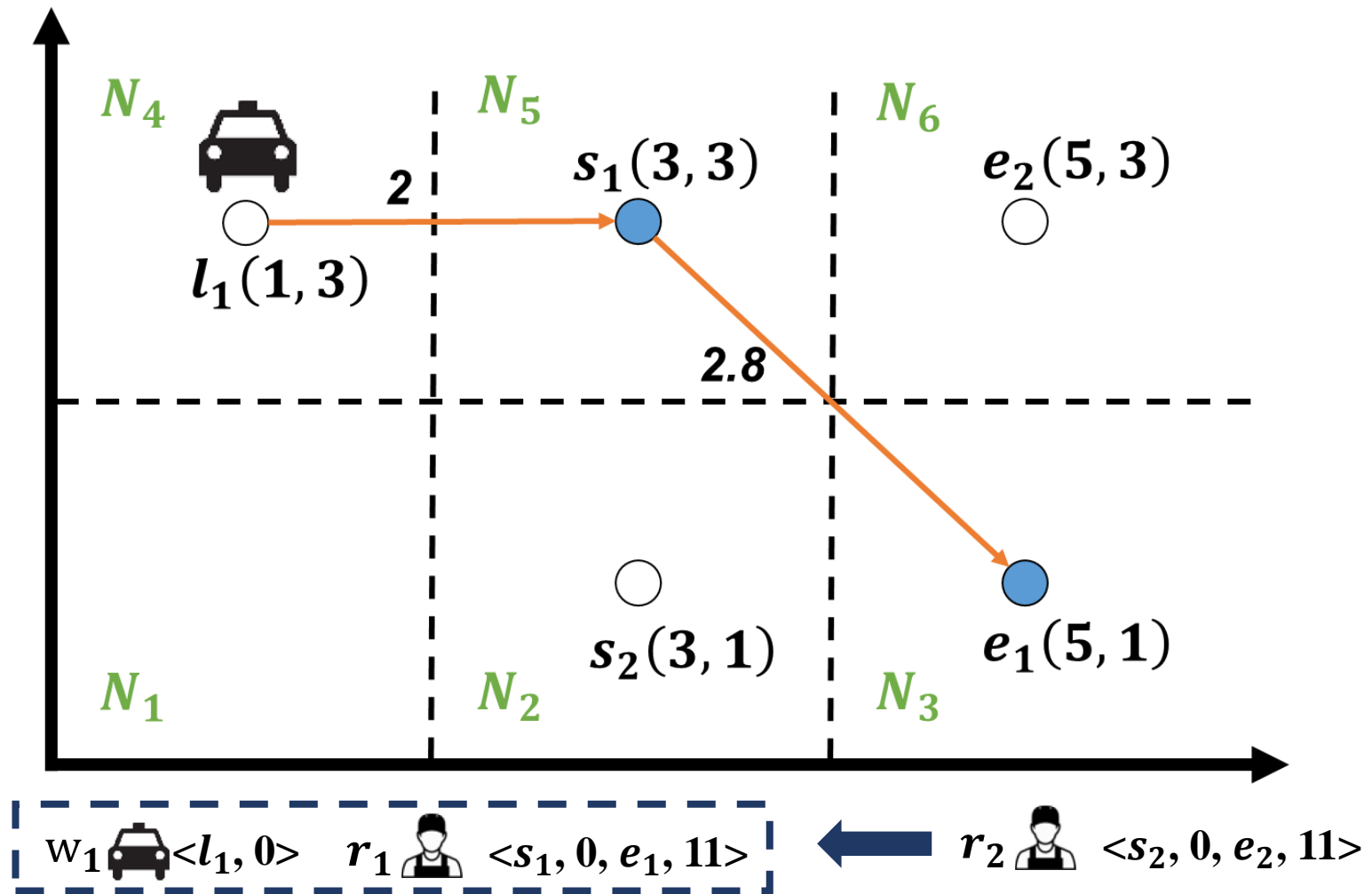
We prove the DARP problem is **NP-hard** by reducing it from the basic route planning problem\* for shareable mobility services. We further show that **no** deterministic nor randomized algorithm can guarantee a **constant Competitive Ratio**



# Running Example

Time spans  $[0 \sim 3, 3 \sim 6, \dots, 15 \sim 18]$

Areas  $[N_1, N_2, \dots, N_6]$



How to derive cost if we assign request  $r_2$ ?



# Running Example

Time spans  $[0\sim 3, 3\sim 6, \dots, 15\sim 18]$   
 Areas  $[N_1, N_2, \dots, N_6]$

**Table 1: Supply Number Map**

$\mathcal{T} \backslash \mathcal{N}$	$N_1$	$N_2$	$N_3$	$N_4$	$N_5$	$N_6$
$T_1$	1.7	3.8	2.5	2.3	0.5	1.3
$T_2$	3.3	2.1	1.7	1.1	3.2	2.9
$T_3$	3.5	3.3	2.0	0.7	3.8	1.4
$T_4$	3.6	1.3	2.4	3.0	1.2	2.6
$T_5$	0.5	2.5	1.4	1.3	1.6	2.3
$T_6$	3.4	2.0	1.0	3.7	2.2	3.8

**Table 2: Demand Number Map**

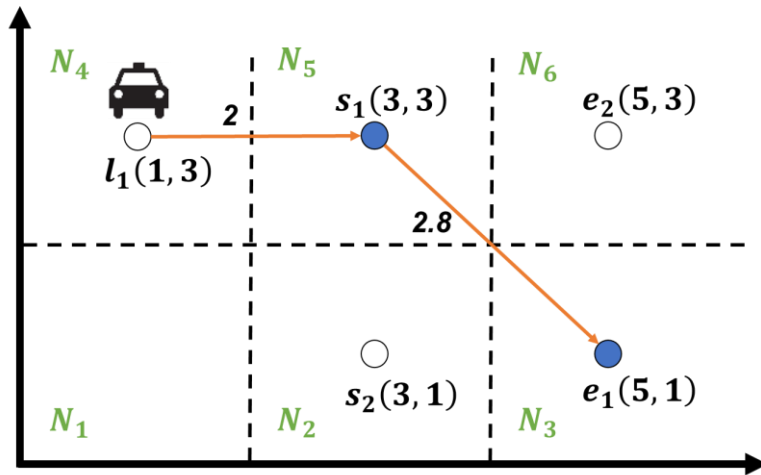
$\mathcal{T} \backslash \mathcal{N}$	$N_1$	$N_2$	$N_3$	$N_4$	$N_5$	$N_6$
$T_1$	2	3	5	2	4	3
$T_2$	3	2	4	2	3	2
$T_3$	3	3	4	2	2	2
$T_4$	3	1	4	3	2	2
$T_5$	4	2	4	3	1	3
$T_6$	3	2	3	4	2	2

Supply number map (**SN**) and Demand number map (**DN**) are required for cost of Demand-Supply Balance (**DSB**)

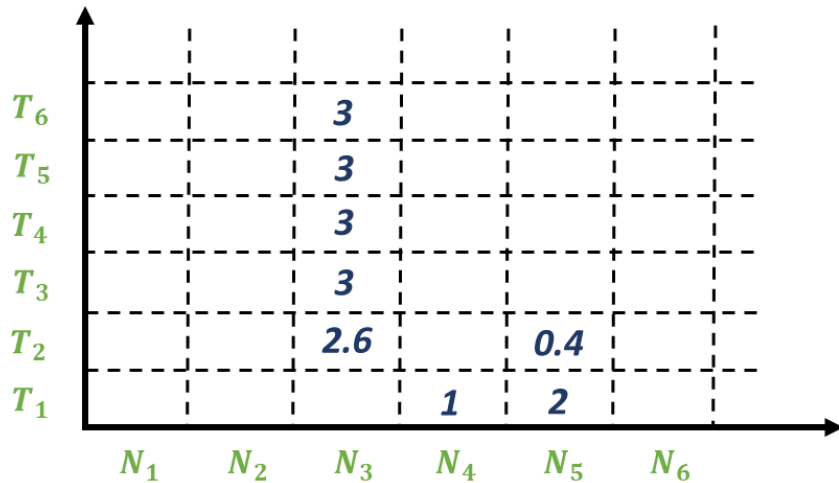


# Running Example

Route before insertion



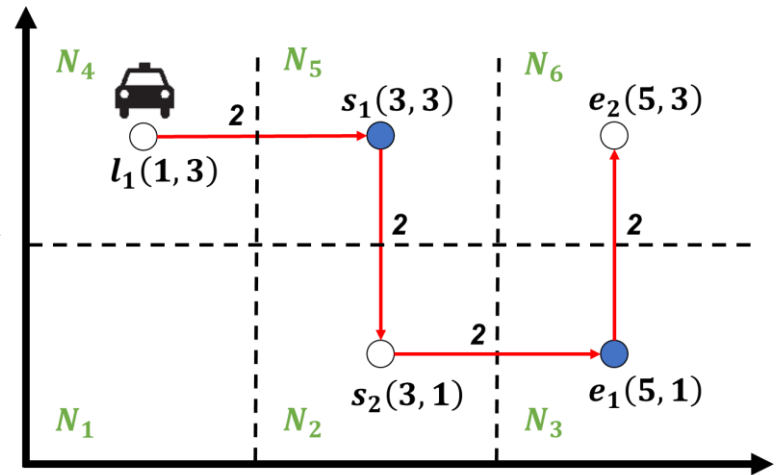
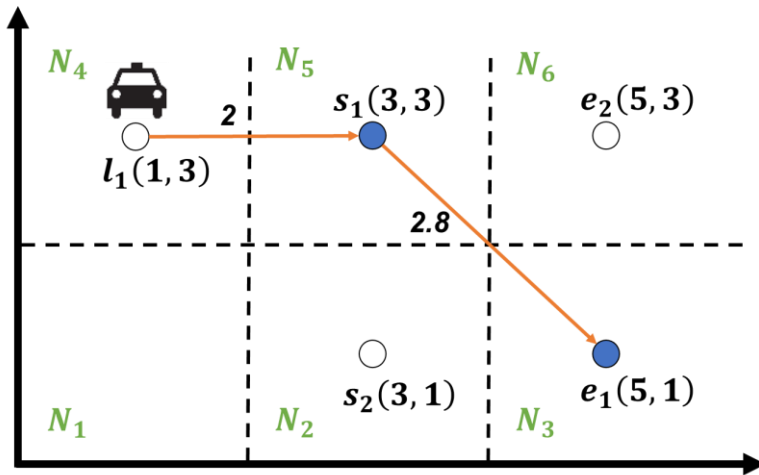
Time duration of  $w_i$  in spatiotemporal cells of SN before insertion



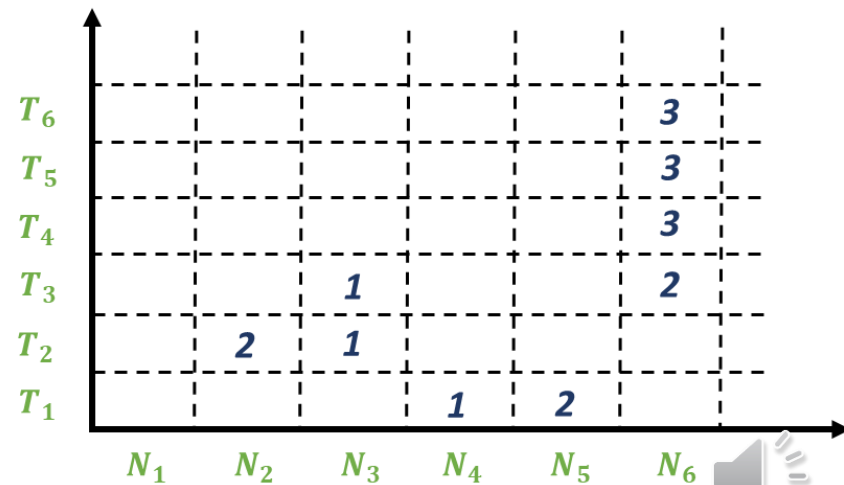
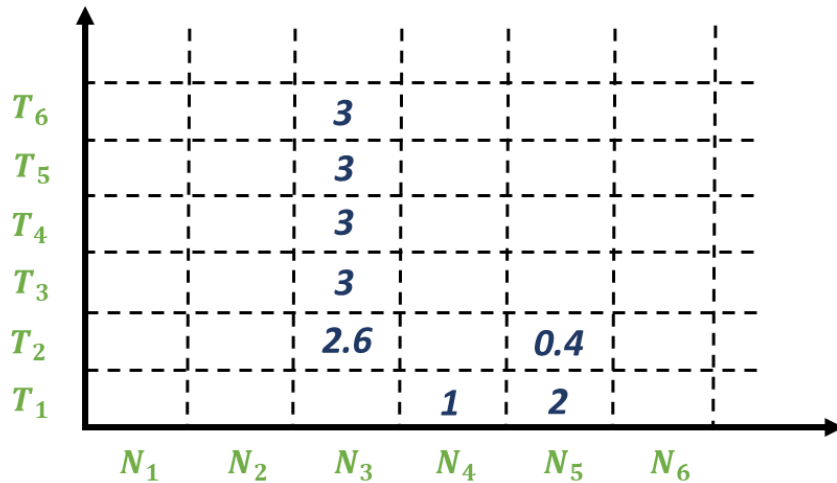
# Running Example

A possible insertion

Route before/after insertion



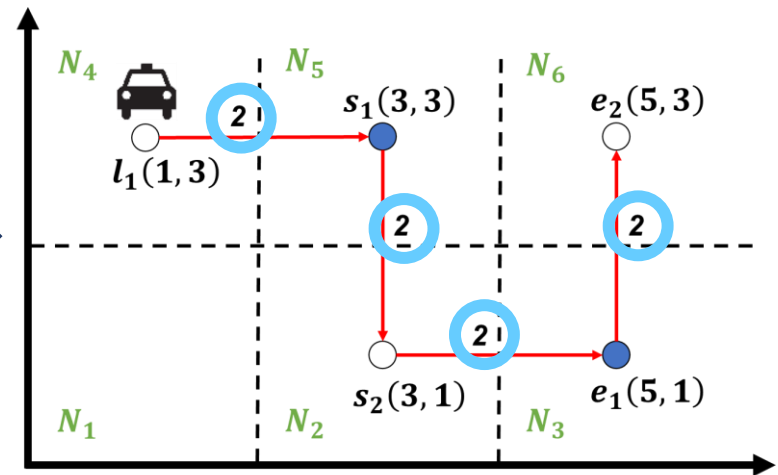
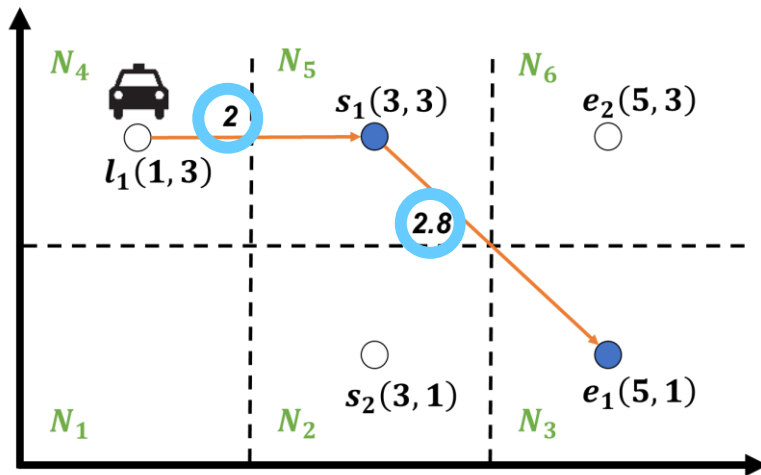
Time duration of  $w_i$  in spatiotemporal cells of SN before/after insertion



# Running Example

A possible insertion

Route before/after insertion



Shortest distance to finish original route  
 $2+2.8$

Shortest distance to finish inserted route  
 $2+2+2+2$



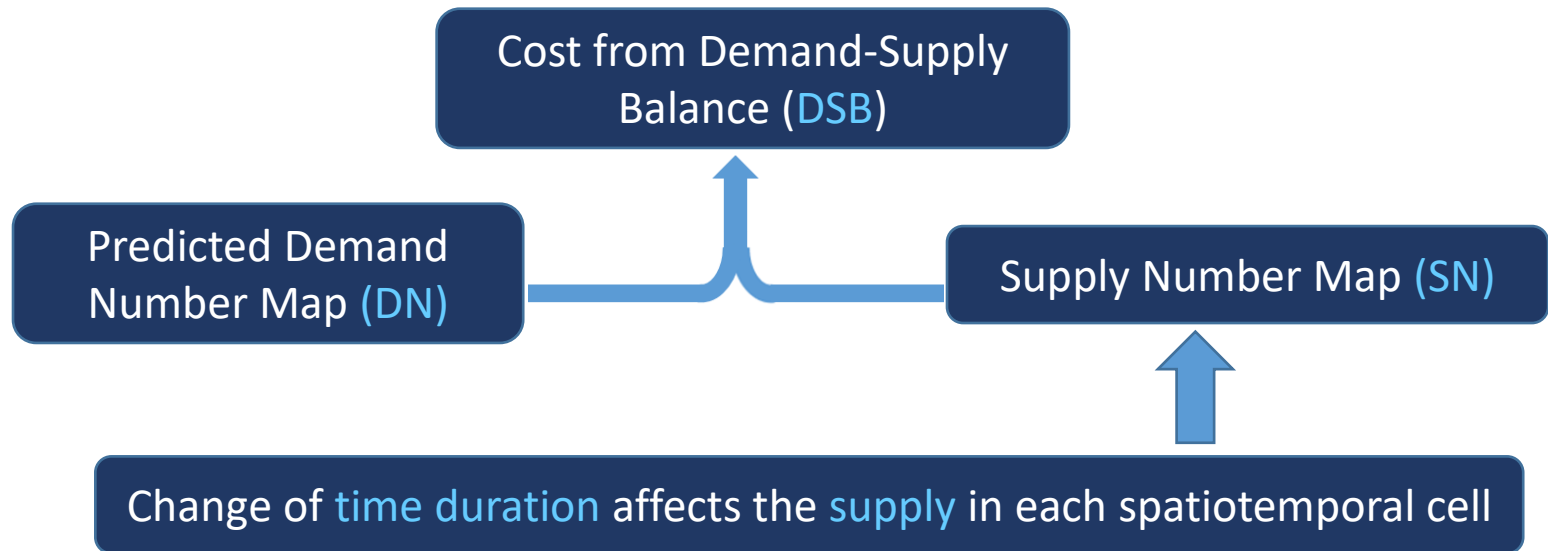
Cost from workers' moving distances  
is derived according to the difference of finishing time



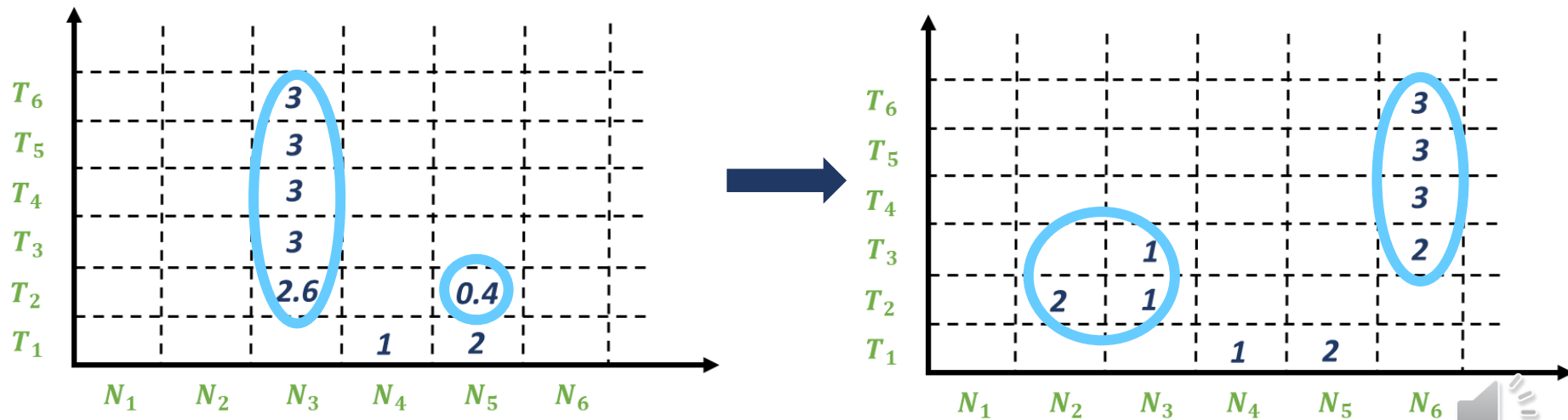


# Running Example

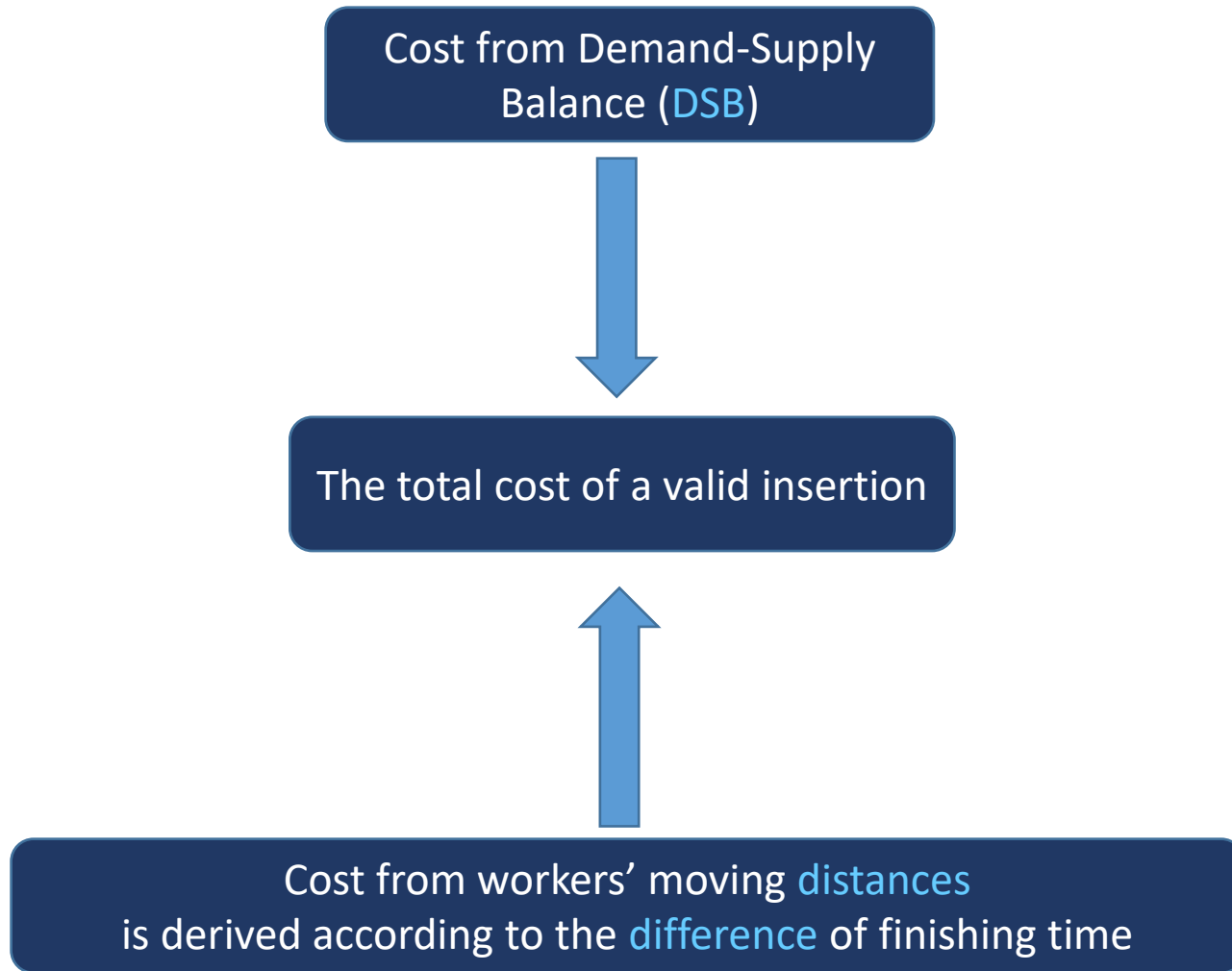
A possible insertion



Time duration of  $w_i$  in spatiotemporal cells of SN before/after insertion



# Running Example



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# Proposed Approaches

To solve the DARP problem, we proposed

## Insertion algorithm (single request)

- **Basic** insertion
- **Dynamic programming**-based insertion

## Solution for DARP problem

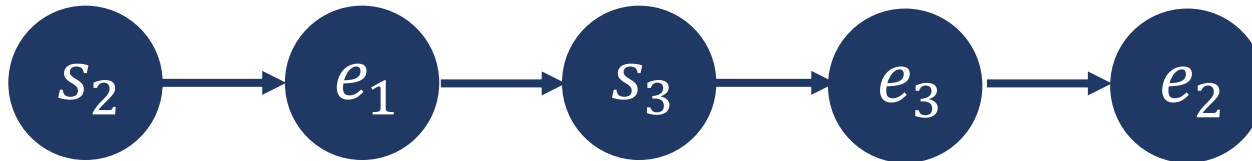
- Insertion-based **dual-phase framework**



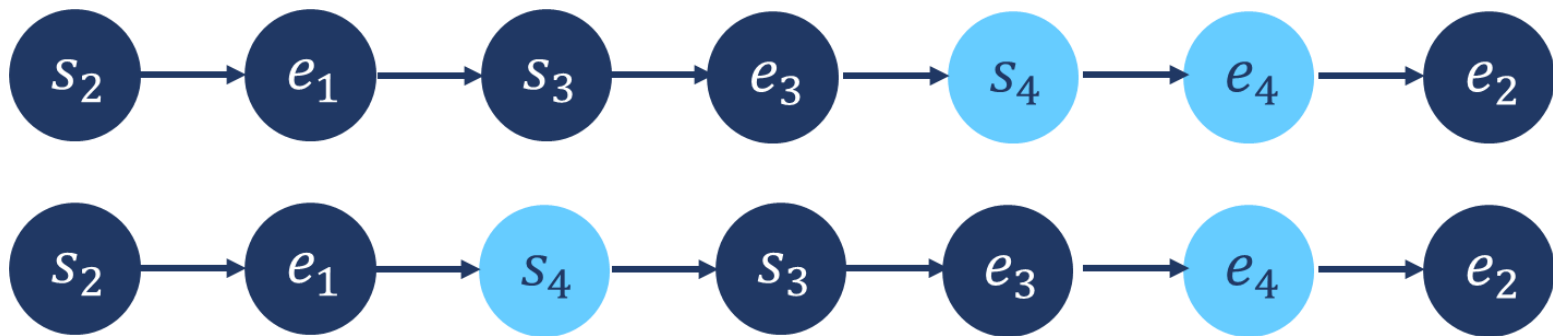
# Insertion

**Insertion:** one request  $\rightarrow$  one worker's route

effective and efficient approach



**Original nodes are in the same order** (search space  $\downarrow$ ):



# Insertion

Naturally:  $O(N^3)$  time complexity

**Distance**-related cost:  $O(N)$ .

Existing work\* reduces its cost as: additional distances from inserting source and destination are **separable**.

**Demand supply balance** cost: previous  $O(N^2)$  and  $O(N)$  algorithms are **not** applicable

Detour of inserting source affects all the supply from latter nodes.



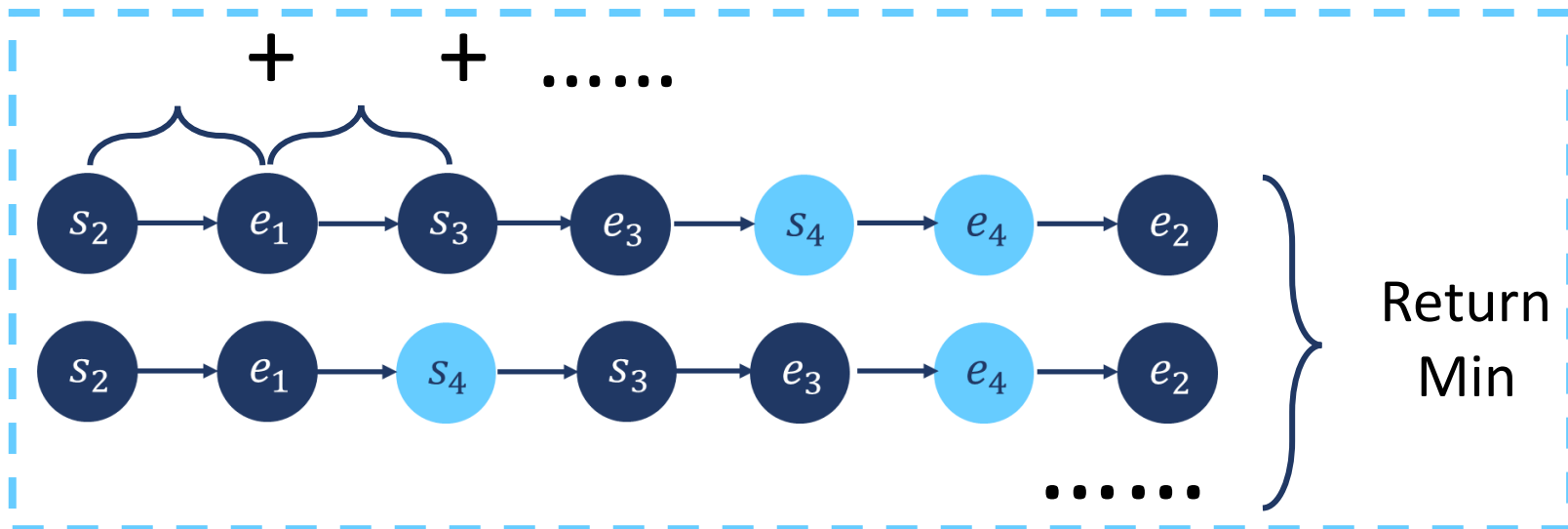
# The Basic Insertion Algorithm ( $O(N^3)$ )

1. **Enumerate** insertion pairs for a length-N route

➤  $O(N^2)$  cases

2. For **each** new route, derive the cost

➤  $N + 1$  small paths. **Calculate** and **sum up** them cost  $O(N)$



3. Return the plan with **lowest** cost (Greedy)



# The DP-Based Insertion Algorithm ( $O(N^2)$ )

## 1. Enumerate insertion pairs

- $O(N^2)$  cases

## 2. For each new route, derive the cost in $O(1)$ time

- **Distance**-related cost in  $O(1)$ : studied\*
- Cost from Demand Supply Balance (**DSB**): how?





# The DP-Based Insertion Algorithm ( $O(N^2)$ )

1. Dynamically derive a **check-up table** in  $O(N)$  time

- Derive the maximum time to delay for each node.
- Divide it into a discretized space. Increasing DSB is stored with **time delay**.

2. Enumerate insertion pairs

- $O(N^2)$  cases

3. For each new route, derive the cost **in  $O(1)$  time**

- Distance-related cost in  $O(1)$ : studied\*
- Cost from DSB: **check in  $O(1)$  time according to time delay**

4. Return the plan with lowest cost



# The DAIF framework

- Assign requests **one-by-one**
- Quickly derive a **lower bound** of cost for each worker
  - In  $O(N)$  time + only **1** shortest path query.
  - Existing work\* derive the lower bound for **distance cost**
  - We efficiently derive a lower bound for **balance cost**  
based on the property of Demand-Supply-Balance cost
- **Sort** → Calculate **exact** cost → **Prune & insert**
  - Derive **exact** cost for each worker ordered by lower bound
  - If the lower bound is larger than current minimal cost, **safely prune** all the rest workers



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# Experimental Setting

- Road Network
  - NYC ( $|V|=61,298$ ,  $|E|=141,372$ )
- Real Dataset
  - Taxi Trips (2013) in NYC (427,093 trip records)
- Synthetic Dataset
  - Generated according to the distribution of NYC (452,116 trip records)



# Experimental Setting

- Compared parameters
  - $e_r$ : the deadline coefficient.
  - $a_i$ : the capacity of workers.
  - $\alpha, \beta$ : the weight for distance/balance cost.
  - $\gamma$ : the factor that staying time duration of worker transfer to supply.
  - $p_o$ : the ratio of penalty cost
  - $|W|$ : number of workers
  - $g$ : grid size

Parameters	Settings
Deadline Coefficient $e_r$	0.1, 0.2, <b>0.3</b> , 0.4, 0.5
Capacity $a_i$	2, <b>3</b> , 4, 7, 10
Distance Weight $\alpha$	<b>1</b>
Balance Weight $\beta$	$\left[ p_r^*, \frac{p_r^*}{e}, \frac{p_r^*}{e^2}, \dots, \frac{p_r^*}{e^5} \right]$
Supply Coefficient $\gamma$	<b>0.0016</b>
Penalty ratio $p_o$	<b>30</b>
Number of workers $ W $	500, 1k, <b>3k</b> , 5k, 10k
Grid size $g$	1k $\times$ 1k, <b>2k <math>\times</math> 2k</b> , 4k $\times$ 4k



# Experimental Setting

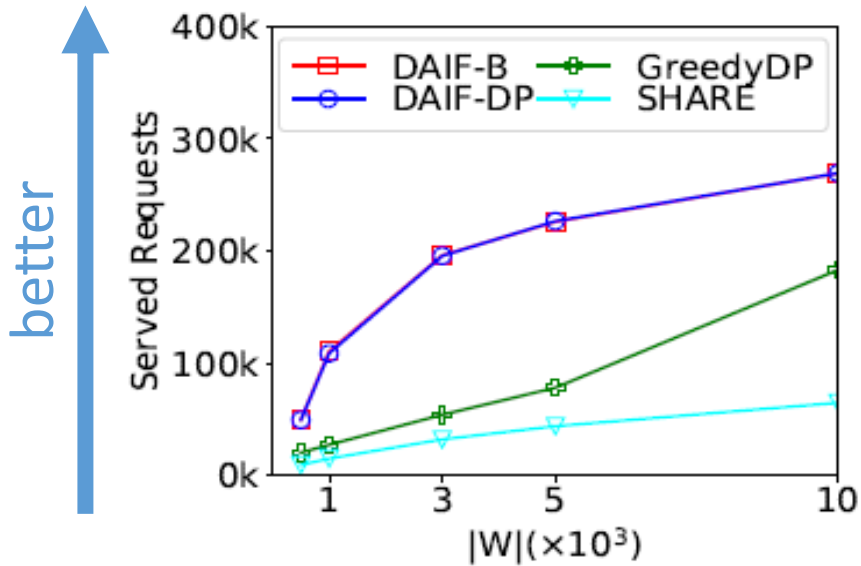
- Tested Algorithms
  - **GreedyDP\***: the state-of-art route planning algorithm using insertion. No demand-related information is used.
  - **SHARE#**: It uses historical information of nodes to choose a route with a higher possibility to pick up passengers along the route
  - **DAIF-B**: our DAIF framework using Basic insertion
  - **DAIF-DP**: our DAIF framework using DP-based insertion

\*Yongxin Tong, et al.: A Unified Approach to Route Planning for Shared Mobility. *Proc. VLDB Endow.* 11(11): 1633-1646 (2018)

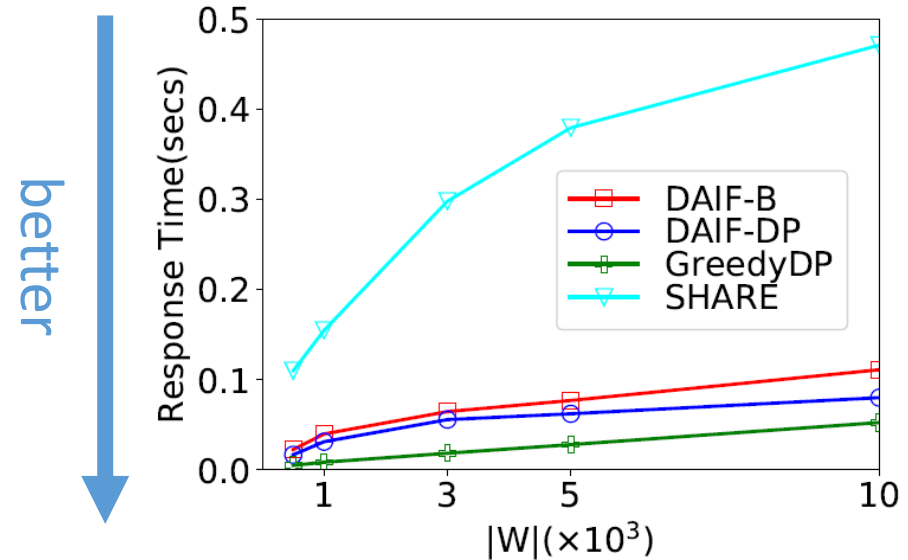
#Chak Fai Yuen, et al.: Beyond Shortest Paths: Route Recommendations for Ride-sharing. *WWW 2019*: 2258-2269



# Experimental Results



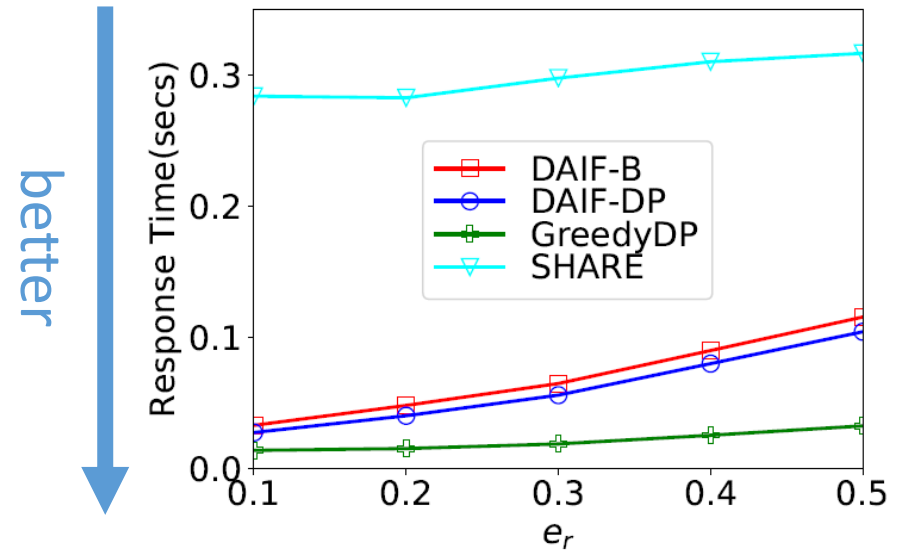
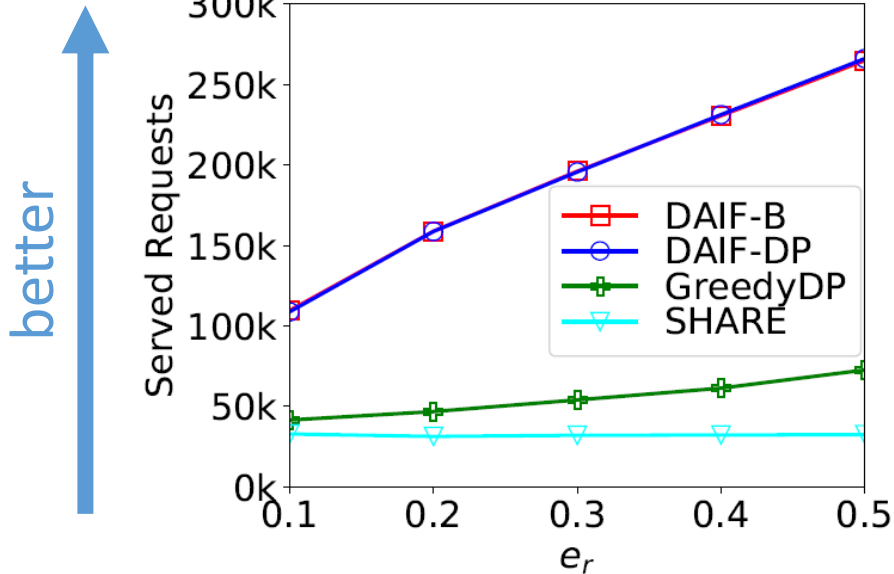
Serve **47%** to **624%**  
more requests



Response time < **0.15s**

Performance of varying number of workers  $|W|$

# Experimental Results



Serve **161.7%** to **718.1%**  
more requests

Response time < **0.15s**

Performance of varying deadline coefficient  $e_r$





# Thank You

## Q&A

The code and datasets  
<https://github.com/dominatorX/DAIF>

